

Introduction

In quantum computation and quantum information transmission, quantum bits (qubits) are susceptible to noise, which can alter their quantum states. Therefore, quantum error-correcting codes are essential for protecting quantum information against noise, while being carefully designed to avoid collapsing the original quantum state and thereby destroying the encoded information.

Motivation: 1-Mode Rotation Codes

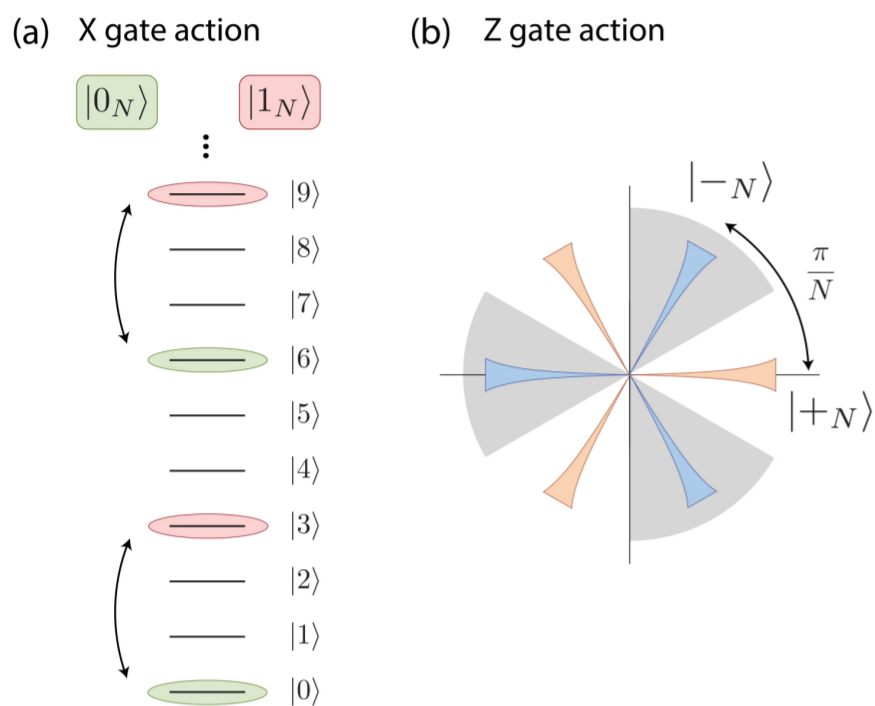


Figure 1: Action of gates on order-3 1-mode rotation code in Fock space (a) and phase space (b).

The 1-mode bosonic rotation codes protect quantum information using rotational symmetry in phase space and inspire multi-mode generalizations such as Tartan code and Chessboard code.

- The computational basis code words are

$$|0_N\rangle = \sum_{k=0}^{\infty} f_{2kN} |2kN\rangle, \\ |1_N\rangle = \sum_{k=0}^{\infty} f_{(2k+1)N} |(2k+1)N\rangle$$

- The dual basis code words are

$$|\pm_N\rangle = \frac{1}{\sqrt{2}} (|0\rangle \pm |1\rangle)$$

- The Z and X stabilizers are

$$\hat{S}_Z = \exp\left[i\frac{2\pi}{N}\hat{n}\right] \text{ and } \hat{S}_X = \sum_{n=0}^{\infty} |n\rangle \langle n+2N|$$

- The logical \bar{Z}_N and \bar{X}_N gates are

$$\bar{Z}_N = \exp\left[i\frac{\pi}{N}\hat{n}\right] \text{ and } \bar{X}_N = \sum_{n=0}^{\infty} |n\rangle \langle n+N|$$

- The logical \overline{CZ} gate for two modes with orders N_1, N_2 is

$$\overline{CZ} = \exp\left[i\frac{\pi}{N_1 N_2} \hat{n}_1 \otimes \hat{n}_2\right]$$

- Errors are described by the basis

$$\hat{E}_m(\theta) = e^{i\theta\hat{n}} \hat{E}_m \text{ where } \hat{E}_m = \sum_{n=0}^{\infty} |n\rangle \langle n+m|.$$

Here \hat{E}_m lowers Fock states by m without \sqrt{n} factors, modelling number-shift errors in rotation codes.

- Correctable errors for an ideal order- N 1-mode code satisfy

$$m \in \{0, \dots, N-1\} \text{ (number errors)}, \\ \theta \in \left(-\frac{\pi}{2N}, \frac{\pi}{2N}\right) \text{ (phase errors)}$$

- The number and rotation error distance of an ideal order- N 1-mode code are

$$d_n = N \text{ (number distance)}, \\ d_\theta = \pi/N \text{ (phase distance)}$$

- The number-phase error trade-off is

$$d_n d_\theta = \pi$$

2-Mode Code: Tartan Codes

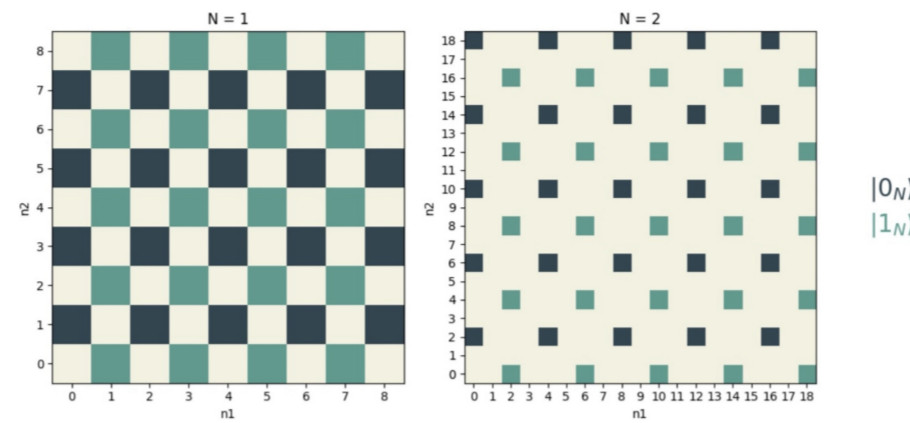


Figure 2: Tartan codes structure for 2-mode encoding with order $N = 1$ and 2. The light and dark green patterns encode the logical states $|1_N\rangle$ and $|0_N\rangle$ respectively (credit: M.Murali).

- The computational basis code words are

$$|0_N\rangle = \sum_{k=0}^{\infty} \sum_{l=0}^{\infty} g_{2k,2l+1} |2Nk\rangle \otimes |N(2l+1)\rangle, \\ |1_N\rangle = \sum_{k=0}^{\infty} \sum_{l=0}^{\infty} g_{2k+1,2l} |N(2k+1)\rangle \otimes |2Nl\rangle$$

- The dual basis code words are defined similarly

$$|\pm_N\rangle = \frac{1}{\sqrt{2}} (|0\rangle \pm |1\rangle)$$

- The Z and X Stabilizers are

$$\hat{S}_{Z_1} = \exp\left[i\frac{2\pi}{N}\hat{n}_1\right] \text{ and } \hat{S}_{Z_2} = \exp\left[i\frac{2\pi}{N}\hat{n}_2\right], \\ \hat{S}_{X_1} = \hat{E}_{2N} \otimes \mathbb{1} \text{ and } \hat{S}_{X_2} = \mathbb{1} \otimes \hat{E}_{2N}$$

where $\hat{n}_1 \equiv \hat{n} \otimes \mathbb{1}$ and $\hat{n}_2 \equiv \mathbb{1} \otimes \hat{n}$

- The logical \bar{Z}_N and \bar{X}_N gate are

$$\bar{Z}_N = \exp\left[i\frac{\pi}{N}\hat{n}_1\right] \text{ and } \bar{X}_N = \hat{E}_N \otimes \hat{E}_N$$

- The logical $\overline{Controlled-Z}$ gate is

$$\overline{CZ} = \exp\left[i\frac{\pi}{N^2} \hat{n}_1 \otimes \hat{n}_2\right]$$

- Both modes can have errors $\hat{E}_m(\theta)$

- Looking at the action of stabilizers on a generic error state

$$|\psi_E\rangle = \left(\hat{E}_{m_1}(\theta_1) \otimes \hat{E}_{m_2}(\theta_2)\right) (\alpha |0_N\rangle + \beta |1_N\rangle)$$

we see

$$\hat{S}_{Z_1} \hat{S}_{Z_2} \hat{S}_{X_1} \hat{S}_{X_2} |\psi_E\rangle \\ \approx \exp\left[\frac{2\pi i}{N} m_1 + \frac{2\pi i}{N} m_2 + i2N\theta_1 + i2N\theta_2\right] |\psi_E\rangle$$

and the syndrome indicates the set of errors $\{\hat{E}_{m_1}(\theta_1) \otimes \hat{E}_{m_2}(\theta_2)\}$ are correctable for order- N Tartan code when

$$m_1, m_2 \in \{0, \dots, N-1\}, \\ \theta_1, \theta_2 \in \left(-\frac{\pi}{2N}, \frac{\pi}{2N}\right)$$

- The conjectured number and rotation error distance of an order- N Tartan code are

$$d_{m_1} = d_{m_2} = N, \\ d_{\theta_1} = d_{\theta_2} = \pi/N$$

- The number-phase error trade-off is

$$(d_{m_1} + d_{m_2})(d_{\theta_1} + d_{\theta_2}) = 4\pi$$

2-Mode Code: Chessboard Codes

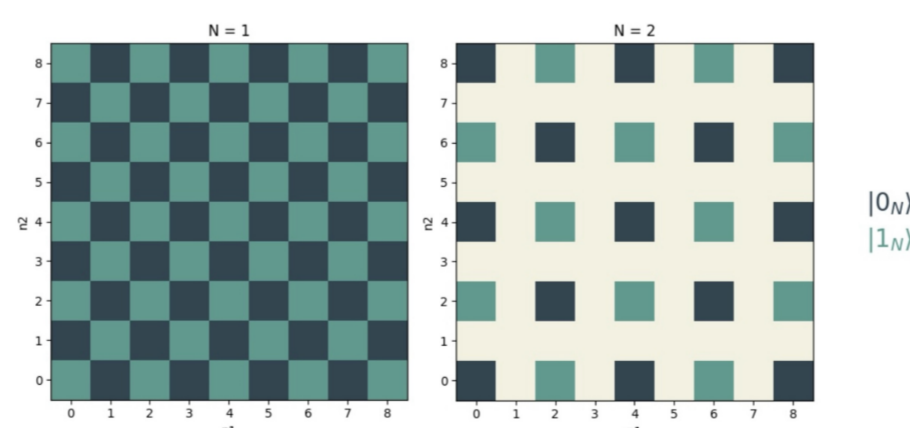


Figure 3: Chessboard codes structure for 2-mode encoding with order $N = 1$ and 2. The light and dark green patterns encode the logical states $|1_N\rangle$ and $|0_N\rangle$ respectively (credit: M.Murali).

- The computational basis code words are

$$|0_N\rangle = \sum_{k=0}^{\infty} \sum_{n_1=0}^{\infty} C_{n_1, n_2} \delta_{n_1, 2jN} \delta_{n_1 + n_2, 2kN} |n_1\rangle \otimes |n_2\rangle, \\ |1_N\rangle = \sum_{k=0}^{\infty} \sum_{n_1=0}^{\infty} C_{n_1, n_2} \delta_{n_1, 2jN} \delta_{n_1 + n_2, (2k+1)N} |n_1\rangle \otimes |n_2\rangle$$

- The dual basis code words are defined similarly

$$|\pm_N\rangle = \frac{1}{\sqrt{2}} (|0\rangle \pm |1\rangle)$$

- The Z and X Stabilizers are

$$\hat{S}_{Z_{\pm}} = \exp\left[i\frac{2\pi}{N}(\hat{n}_1 \pm \hat{n}_2)\right], \\ \hat{S}_{X_{p,q}} = \hat{E}_p \otimes \hat{E}_q$$

where $p, q \in \mathbb{Z}_{\geq 0}$ and $(p+q = 2N$ or $|p-q| = 2N)$

- The logical \bar{Z}_N and \bar{X}_N gate are

$$\bar{Z}_{N_{\pm}} = \exp\left[i\frac{\pi}{N}(\hat{n}_1 \pm \hat{n}_2)\right], \\ \bar{X}_N = \hat{E}_p \otimes \hat{E}_q$$

where $p, q \in \mathbb{Z}_{\geq 0}$ and $(p+q = N$ or $|p-q| = N)$

- The logical $\overline{Controlled-Z}$ gate is

$$\overline{CZ}_{N_1 N_2} = \exp\left[i\frac{\pi}{N_1 N_2}(\hat{n}_1 \pm \hat{n}_2) \otimes (\hat{n}_1 \pm \hat{n}_2)\right]$$

- Looking at the action of stabilizers on a generic error state

$$|\psi_E\rangle = \left(\hat{E}_{m_1}(\theta_1) \otimes \hat{E}_{m_2}(\theta_2)\right) (\alpha |0_N\rangle + \beta |1_N\rangle)$$

we see

$$\hat{S}_{Z_+} \hat{S}_{Z_-} \hat{S}_{X_{2N,0}} \hat{S}_{X_{0,2N}} |\psi_E\rangle \\ \approx \exp\left[\frac{2\pi i}{N}(m_1 + m_2)\right] \exp\left[\frac{2\pi i}{N}(m_1 - m_2)\right] \\ \exp[i2N\theta_1 + i2N\theta_2] |\psi_E\rangle$$

and the set of errors $\{\hat{E}_{m_1}(\theta_1) \otimes \hat{E}_{m_2}(\theta_2)\}$ are correctable for order- N Chessboard code when

$$m_1 + m_2 \in \{0, \dots, N-1\}, \\ m_1 - m_2 \in \{0, \dots, N-1\}, \\ \theta_1, \theta_2 \in \left(-\frac{\pi}{2N}, \frac{\pi}{2N}\right)$$

- The conjectured number and rotation error distance of an order- N Chessboard code are

$$d_{m_1} = d_{m_2} = N, \\ d_{\theta_1} = d_{\theta_2} = \pi/N$$

- The number-phase error trade-off is

$$(d_{m_1} + d_{m_2})(d_{\theta_1} + d_{\theta_2}) = 4\pi$$

Discussion and Conclusion

The number and rotation error distances of both the Tartan and Chessboard codes are conjectured based on the syndromes of the corresponding error states. A natural direction for future work is to employ the Knill-Laflamme conditions to prove the conjectured distances.

Both two-mode codes enhance the number-phase error trade-off from π to 4π due to the additional mode, indicating that the logical qubits constructed from these codes are more tolerant to noise during quantum information transmission.

Acknowledgement

I express my gratitude to my supervisor Dr. Joshua Combes for his support; to my friend and colleague Mugilan Murali for the valuable discussions which inspired me deeply, and the outstanding figures he made for this poster; and to the School of Mathematics and Statistics for providing me with this research opportunity.

References

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